

A Very Strong Higher Guessing Model Principle

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IM PAN



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- ▶ *Martin's Axiom*: c.c.c. forcings and κ arbitrary.
- ▶ Increasing κ or enlarging \mathcal{C} yields stronger axioms.

$$\text{MM} \Leftrightarrow \text{SPFA} \Rightarrow \text{PFA} \Rightarrow \text{MA}_{\omega_1}$$

Old but gold

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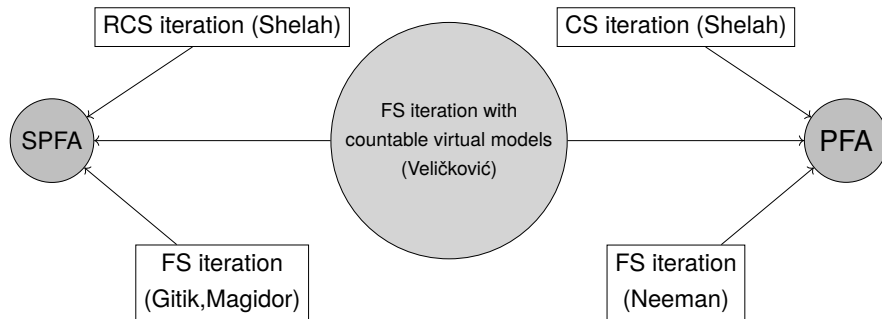
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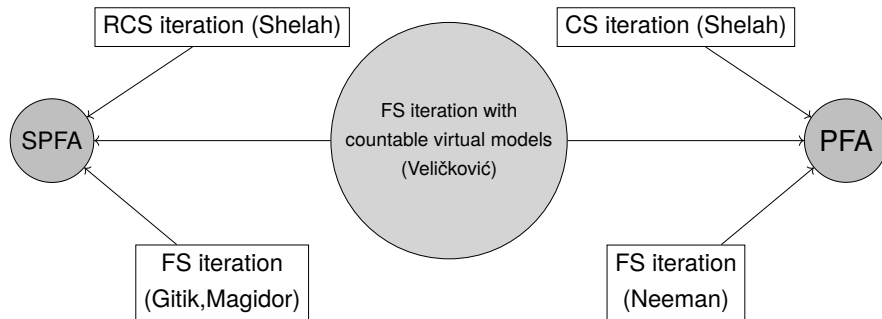
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Do we have forcing axioms implying certain statements like $|\mathbb{R}| \geq \omega_3$, $\text{TP}(\omega_3)$, SCH, etc.?

Iteration and forcing axioms



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No iteration theorem for stationary-set preserving forcings.

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Basically, it’s the axiom we can trust to have our back.

I will discuss certain consistency results that could be regarded as potential consequences of an imaginary higher forcing axiom.

Guessing models

Definition (Viale, Weiss)

$M \prec \mathcal{H}(\theta)$ is a κ -guessing model if for every $x \in \bigcup_{X \in M} \mathcal{P}(X)$ the following are equivalent.

1. $\exists x^* \in M$ so that $x \cap M = x^* \cap M$,
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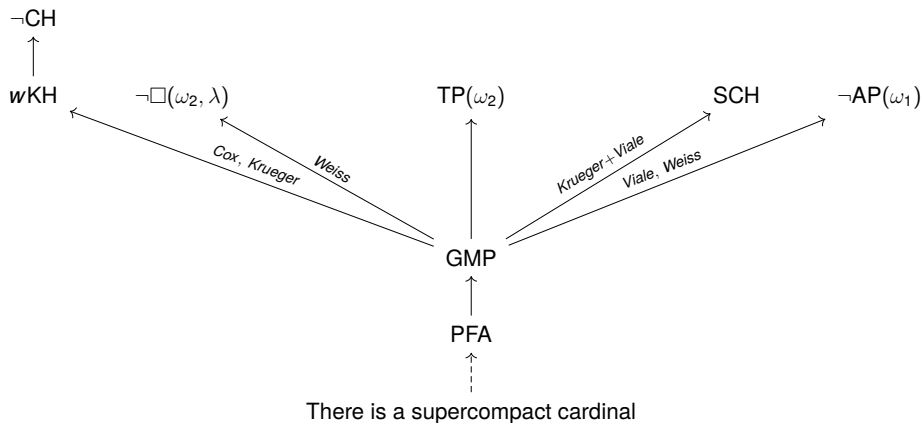
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Theorem (Viale, Weiss)

PFA implies GMP: $\{M \prec \mathcal{H}(\theta) : |M| = \omega_1 \wedge M \text{ is } \omega_1\text{-guessing}\}$ is stationary for every $\theta \geq \omega_2$.

Consequences



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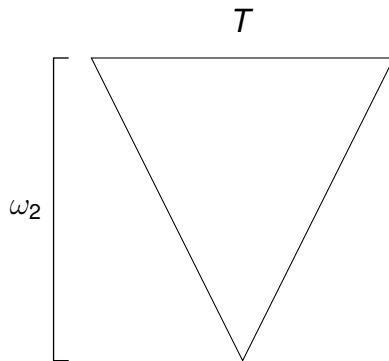
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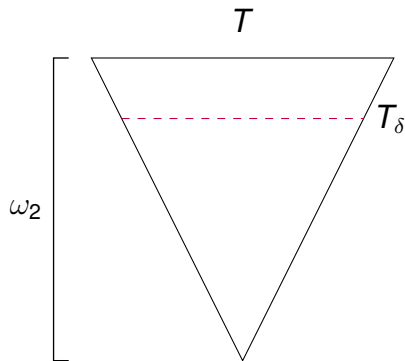
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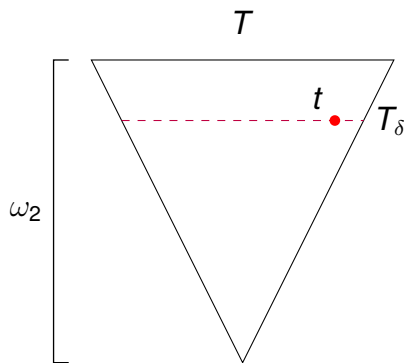
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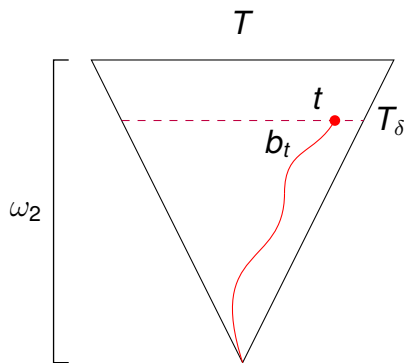
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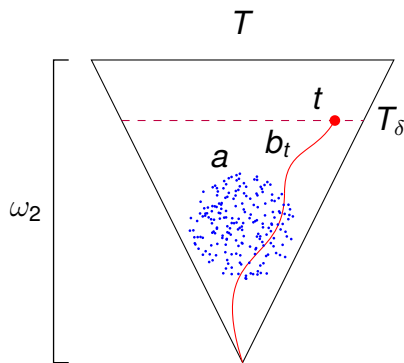
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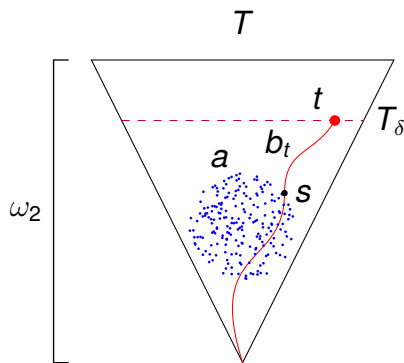
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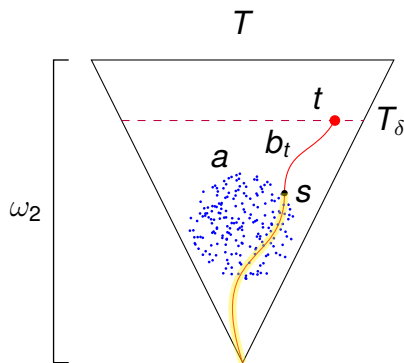
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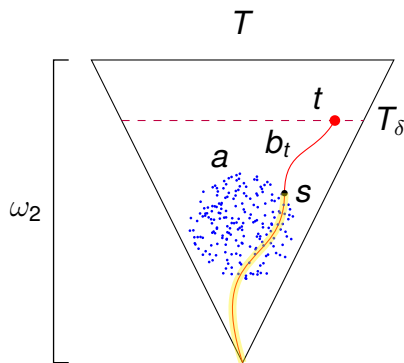
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$$a \cap b_t = a \cap b_s \in M$$

Approachability

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$\delta < \omega_2$ is *approachable by \vec{a}* if there is $c \subseteq \delta$ such that

1. $\text{ot}(c) < \delta$.
2. $\forall \xi < \delta \exists \zeta < \delta \ a_\xi \cap c = a_\zeta$.

Approachability ideal

The *approachability ideal* $I[\omega_2]$ is generated by $\mathfrak{J}_{\text{ns}}(\omega_2)$ and sets of the form

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Definition

$\text{AP}(\omega_1)$ states that $\omega_2 \in I[\omega_2]$.

$$\text{GMP} \Rightarrow \neg \text{AP}(\omega_1)$$

Proof. There are stationarily many non-approachable points of cofinality ω_1 .

Fix an approaching sequence $\vec{b} = \langle b_\xi : \xi < \omega_2 \rangle$.

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Proof. There are stationarily many non-approachable points of cofinality ω_1 . Fix an approaching sequence $\vec{b} = \langle b_\xi : \xi < \omega_2 \rangle$. Let $M \prec \mathcal{H}(\theta)$ be an ω_1 -guessing model of size ω_1 with $\vec{b} \in M$.

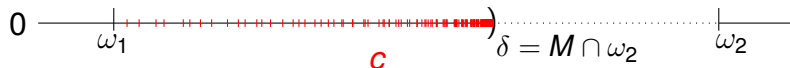
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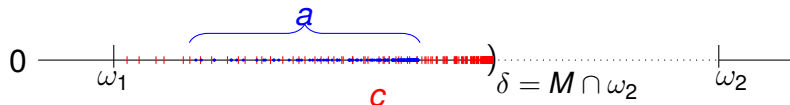
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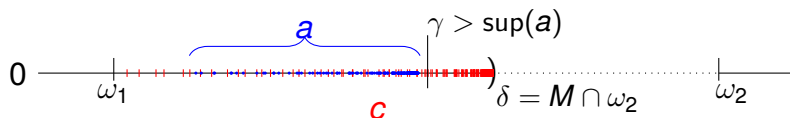
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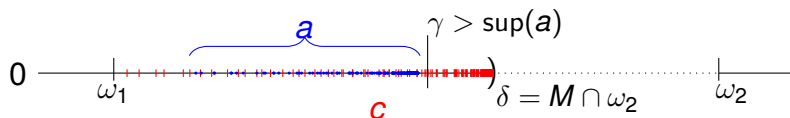
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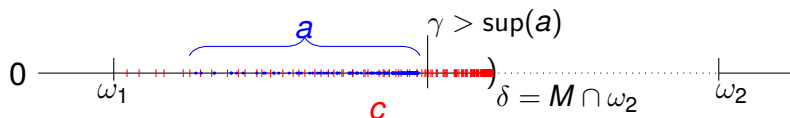
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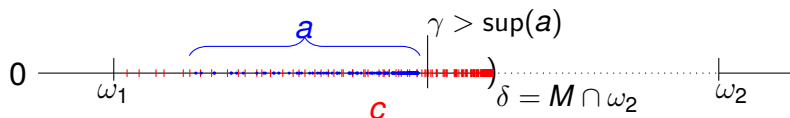


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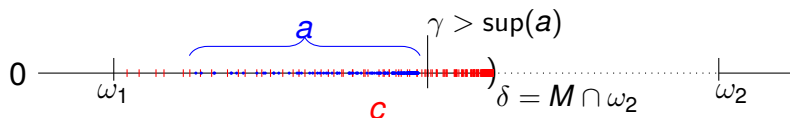


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Let $T := \text{ZFC}$ together with the statement that for every $\theta \geq \omega_3$, the set

$$\{M \prec \mathcal{H}(\theta) : |M| = \omega_2 \wedge M^\omega \subseteq M \wedge M \text{ is } \omega_2\text{-guessing}\}$$

is stationary in $\mathcal{P}_{\omega_3}(\mathcal{H}(\theta))$.

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Theorem (Trang)

1. $\text{Con}(\text{ZFC} + \exists \text{ s.c. cardinal}) \Rightarrow \text{Con}(T)$.
2. Assume T . There is a transitive model $M \models \text{"AD}_{\mathbb{R}} + \Theta \text{ is regular"}$ with $\mathbb{R} \subseteq M$.

GMP⁺

Definition

$M \prec \mathcal{H}(\theta)$ of size ω_2 is a *strongly* ω_1 -guessing model if it is the union of an ω_1 -closed \in -sequence of ω_1 -guessing models of size ω_1 .

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Definition (GMP⁺)

For every $\theta \geq \omega_3$, there are stationarily many strongly ω_1 -guessing elementary submodels of $\mathcal{H}(\theta)$

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M., Veličković

GMP⁺ implies $I[\omega_2] = \mathfrak{J}_{\text{ns}}(\omega_2) \bmod \text{Cof}(\omega_1)$.

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Theorem (Mitchell)

$I[\omega_2] = \mathfrak{J}_{\text{ns}}(\omega_2) \bmod \text{Cof}(\omega_1)$ is consistent.

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The proof uses forcing with pure side conditions. Let $\kappa < \lambda$ be supercompact cardinals. The conditions are finite sets of countable and κ -Magidor virtual models which satisfy certain requirements with respect to $\{\alpha < \lambda : V_\alpha \prec V_\lambda\}$. The forcing is proper for all models involved and forces $\kappa = \omega_2$ and $\lambda = \omega_3$.

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The proof uses forcing with pure side conditions. Let $\kappa < \lambda$ be supercompact cardinals. The conditions are finite sets of countable and κ -Magidor virtual models which satisfy certain requirements with respect to $\{\alpha < \lambda : V_\alpha \prec V_\lambda\}$. The forcing is proper for all models involved and forces $\kappa = \omega_2$ and $\lambda = \omega_3$.

Note that $I[\omega_2] = \mathfrak{J}_{\text{ns}}(\omega_2) \bmod \text{Cof}(\omega_1)$ implies $|\mathbb{R}| \geq \omega_3$ (due to Shelah.)

GMP⁺

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► $\text{GMP}^+ \Rightarrow |\mathbb{R}| \geq \omega_3.$

Generalization

$\text{GMP}(\kappa, \gamma)$ states that for every sufficiently large regular cardinal θ , the set

$$\mathcal{G}_{\kappa, \gamma}(\mathcal{H}(\theta)) := \{M \prec \mathcal{H}(\theta) : |M| < \kappa \wedge M \text{ is } \gamma\text{-guessing}\}$$

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$\text{GMP}^+(\kappa, \gamma)$ is defined in the obvious way. So $\text{GMP} \equiv \text{GMP}(\omega_2, \omega_1)$ and

$\text{GMP}^+ \equiv \text{GMP}^+(\omega_3, \omega_1)$.

- The appropriate generalizations of the consequences of GMP or GMP^+ follow from their corresponding higher principles.

IGMP

Definition (Cox, Krueger)

A γ -guessing model is called *indestructible* if it remains guessing in any outer transitive universe in which γ is a cardinal.

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Theorem (Cox, Krueger)

$PFA \Rightarrow IGMP$.

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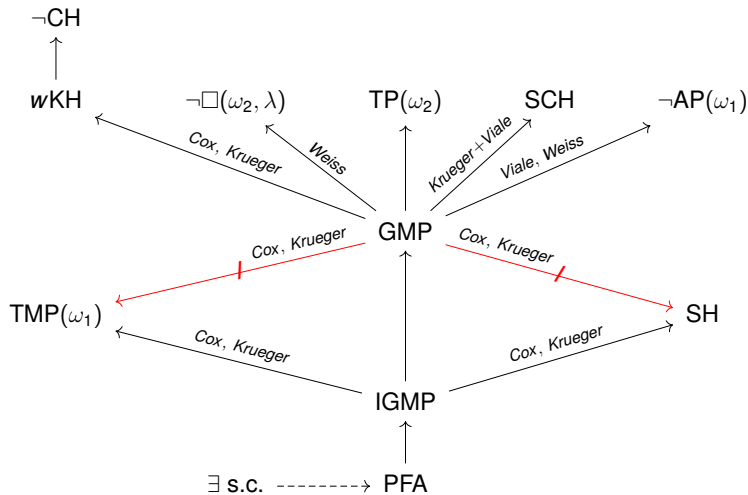
Idea: Iterate standard specializing forcing up to a supercompact cardinal κ using *finite conditions* and then add an arbitrarily large number of reals.
The iteration is technical and delicate.

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Idea: Iterate standard specializing forcing up to a supercompact cardinal κ using *finite conditions* and then add an arbitrarily large number of reals. The iteration is technical and delicate. The main difficulty is performing a quotient analysis for models of size less than κ (the so-called κ -Magidor models).



TMP

Todorčević Maximality Principle.

Definition ($\text{TMP}(\kappa^+)$)

Every forcing that adds a new subset of κ^+ whose initial segments are in the ground model must collapse some cardinal.

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Theorem (Todorćević, $2^{\aleph_0} < \aleph_{\omega_1}$)

Suppose that every tree of size and height ω_1 without cofinal branches is special. Then $\text{TMP}(\omega_1)$ holds.

Theorem (Cox, Krueger, and $2^{\aleph_0} < \aleph_{\omega_1}$)

$IGMP \Rightarrow TMP(\omega_1)$

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Theorem (Golshani, Shelah)

$TMP(\kappa^+)$ is forceable, assuming suitable large cardinals.

IGMP⁺

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Definition

$M \prec \mathcal{H}(\theta)$ of size ω_2 is an *indestructible strongly ω_1 -guessing model* if it is the union of an ω_1 -closed \in -sequence of indestructible ω_1 -guessing models of size ω_1 .

IGMP⁺

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Suppose that $V \subseteq W$ are transitive inner models of ZFC.

► $V \models \text{“IGMP}^+ + 2^{\omega_1} < \aleph_{\omega_2} \text{”}$.

► $\mathcal{P}(\omega_2)^V \neq \mathcal{P}(\omega_2)^W$.

Then either $\mathcal{P}(\omega_1)^V \neq \mathcal{P}(\omega_1)^W$ or some V -cardinal $\leq 2^{\omega_1}$ is no longer a cardinal in W .

IGMP⁺

Corollary ($\text{IGMP}^+ + 2^{\omega_0} < \aleph_{\omega_1} + 2^{\omega_1} < \aleph_{\omega_2}$)

Let $W \models \text{“ZFC”}$ be a transitive extension with the same cardinals and reals.

Then $\mathcal{P}(\omega_2)^V = \mathcal{P}(\omega_2)^W$.

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Corollary ($2^{\aleph_0} < \aleph_{\omega_1} + 2^{\aleph_1} < \aleph_{\omega_2}$)

IGMP⁺ implies both $\text{TMP}(\omega_1)$ and $\text{TMP}(\omega_2)$.

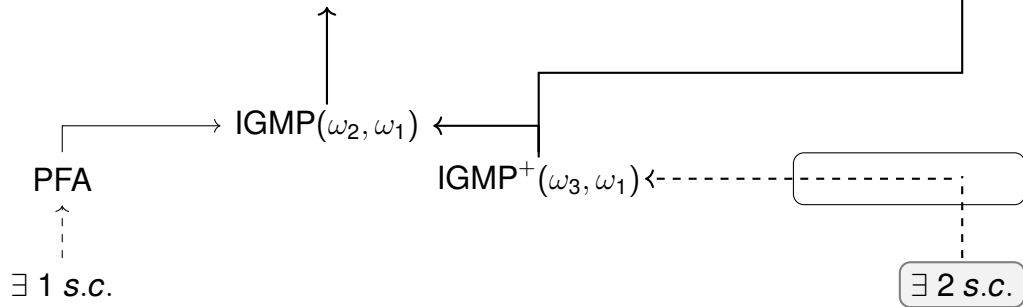
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$$\text{GMP}^+(\omega_3, \omega_1) + \text{TMP}(\omega_2) + \text{TP}(\omega_3) + \neg\Box(\omega_3, \lambda) + I[\omega_2] \sim \mathfrak{J}_{\text{ns}}(\omega_2)$$

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