

# Příkrý Forcing and Properness

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## Abstract

It is shown the Příkrý forcing for measurable  $\kappa$  is proper below  $V_\kappa$ .

Let  $\mathcal{U}$  be a measure on  $\kappa$ . The Příkrý forcing  $\mathbb{P}_\mathcal{U}$  consists of conditions  $p := (A_p, s_p)$ , where

1.  $A_p \in \mathcal{U}$ ,
2.  $s_p \in [\kappa]^{<\omega}$ , and
3.  $\max(s_p) < \min(A_p)$ .

A condition  $p$  is stronger than  $q$ , i.e.  $p \leq q$ , if  $A_p \subseteq A_q$ ,  $s_p \cap \max(s_q) = s_q$ , and  $s_p \setminus s_q \subseteq A_q$ .

**Proposition 0.1.** *For every regular  $\lambda > 2^{\kappa^+}$ , every  $X \in V_\kappa$ , every condition  $p$ , and every  $M \prec H_\lambda$  with  $p, X, \mathcal{U} \in M$  of size  $< \kappa$ , there is  $q \leq p$  such that  $q$  is  $(M, \mathbb{P}_\mathcal{U}, X)$ -generic, i.e.,*

$$q \Vdash_{\mathbb{P}_\mathcal{U}} "M[\dot{G}] \cap X = M \cap X".$$

*Proof.* Let  $p := (A_p, s_p)$ . We first define a condition  $p'$  by letting

$$A_{p'} := \bigcap (M \cap \mathcal{U}),$$

and  $s_{p'} := s_p$ . By  $\kappa$ -completeness of  $\mathcal{U}$ ,  $A_{p'} \in \mathcal{U}$ . Clearly  $p' \leq p$ . Let  $\Omega$  be the set of pairs  $(\dot{\tau}, x)$  such that

- $\dot{\tau}$  is a  $\mathbb{P}_\mathcal{U}$ -name in  $M$ ,
- $x \in X$ , and
- some condition below  $p$  forces  $\dot{\tau} = \check{x}$ .

Clearly  $|\Omega| < \kappa$ . Now by the Příkrý condition<sup>1</sup> and the  $\kappa$ -completeness of  $\mathcal{U}$ , there is  $B \subseteq A_{p'}$  such that  $q := (B, s_p)$  decides  $\dot{\tau} = \check{x}$ , for every  $(\dot{\tau}, x) \in \Omega$ . We claim that  $q$  is as required.

Suppose that  $\dot{\sigma}$  is forced by some condition  $q' \leq q$  to be in  $M[\dot{G}] \cap X$ . We may extend  $q'$  to some condition  $q''$ , find a  $\mathbb{P}_\mathcal{U}$ -name  $\dot{\tau} \in M$ , and an element  $x \in X$  such that  $q''$  forces  $\dot{\tau} = \dot{\sigma} = \check{x}$ . Therefore,  $(\dot{\tau}, x) \in \Omega$ , since  $q'' \leq p$ . On the other hand,  $q$  must decide  $\dot{\tau} = \check{x}$  in the same way, that is  $q \Vdash "\dot{\tau} = \check{x}"$ . Notice that  $q$  is a direct extension of  $p$ . So by elementarity, there are  $A' \in \mathcal{U} \cap M$  and  $x' \in X \cap M$  such that  $(A', s_p)$  forces  $\dot{\tau} = \check{x}'$ . But  $q \leq (A', s_p)$ , and hence  $x' = x$ . Therefore,  $q'' \Vdash \dot{\sigma} \in M \cap X$ . Since  $\dot{\sigma}$  was arbitrary, we have

$$q \Vdash_{\mathbb{P}_\mathcal{U}} "M[\dot{G}] \cap X = M \cap X".$$

□

**Remark 0.2.**  $\mathbb{P}_\mathcal{U}$  preserves stationary sets in  $V_\kappa$ .

**Remark 0.3.** Letting  $X = \omega_1$  and  $M$  be countable, we have that  $\mathbb{P}_\mathcal{U}$  is semi-proper.

<sup>1</sup>For every statement  $\phi$  in the forcing language of  $\mathbb{P}_\mathcal{U}$  and for every  $p \in \mathbb{P}_\mathcal{U}$ , there is  $q \leq p$  with  $s_q = s_p$  that decides  $\phi$ .